

# High Performance Computing Algorithm Design and Novel Computing Environments

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Follow-up Workshop

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  - Minimum Intrusion Grid
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## Target problems

Solve the system of linear equations

$$Ax = b$$

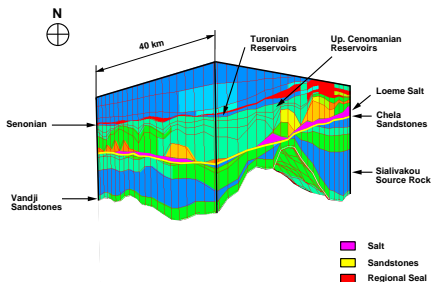
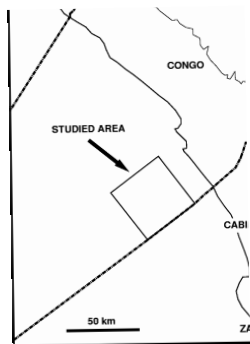
where the matrix  $A$  is:

- sparse,
- large,
- with highly varying coefficients (for example,  $|a_{ij}| \in [10^{-6}, 10^6]$ )



# Lithology example

Sedimentary basin simulation



Position of the studied area

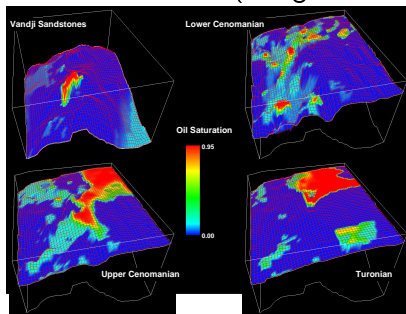
Lithology of the studied 3D-block

(Taken From F.Schneider Et Al. Oil & Gas Science And Technology 55 (1), 2000)



# Congo Offshore oil reservoir simulation

## Numerical Results (Congo Offshore)



Computed oil saturation for the potential reservoir levels at present day

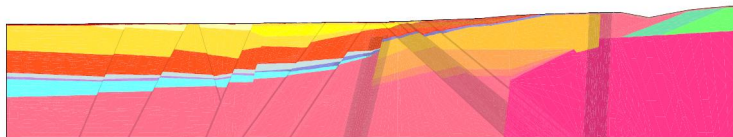
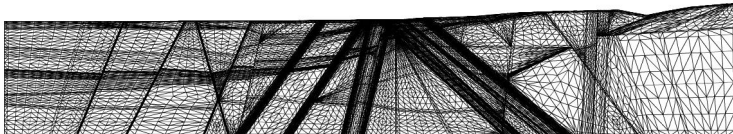
(Taken from {F.Schneider et al. Oil & Gas Science and Technology 55 (1), 2000})



# Sellafield

Groundwater flow, Test case “Sellafield”

Complicated deterministic media ©NIREX UK Ltd.



■	EDZ
■	CROWN SPACE
■	WASTE VAULTS
■	FAULTED GRANITE
■	GRANITE
■	DEEP SKIDDAW
■	N-S SKIDDAW
■	DEEP LATTERBARROW
■	N-S LATTERBARROW
■	FAULTED TOP M-F BVG
■	TOP M-F BVG
■	FAULTED BLEABATH BVG
■	BLEABATH BVG
■	FAULTED F-H BVG
■	F-H BVG
■	FAULTED UNDIFF BVG
■	UNDIFF BVG
■	FAULTED N-S BVG
■	N-S BVG
■	FAULTED CARB LIST
■	CARB LIST
■	FAULTED COLLYHRIST
■	COLLYHRIST
■	FAULTED BROCKRAM
■	BROCKRAM
■	SHALES - EVAP
■	FAULTED BRHM
■	BOTTOM NHM
■	FAULTED DEEP ST BEES
■	DEEP ST BEES
■	FAULTED N-S ST BEES
■	N-S ST BEES
■	FAULTED VN-S ST BEES
■	VN-S ST BEES
■	FAULTED DEEP CALDER
■	DEEP CALDER
■	FAULTED N-S CALDER
■	N-S CALDER
■	FAULTED VN-S CALDER
■	VN-S CALDER
■	MERCIA MUDSTONE
■	QUATERNARY

## Direct methods

UMFPACK, SuperLU, MUMPS, Pardiso

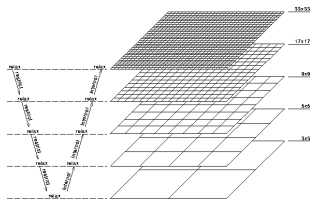
- 1 Analysing the sparsity structure to introduce less fill-in via re-ordering
  - 2 Factorisation step
  - 3 Solving step
- Roughly 100(1)-10(2)-1(3) time factor  
2D - often still OK,  
3D - the first step becomes too expensive, and still, the fill-in starts dominating.



## Iterative methods

Richardson's, Krylov subspace methods, MultiGrid

- Richardson's type iterations
- Krylov subspace methods
- Geometric multigrid
- Algebraic multigrid
  - f-c colouring
  - aggregation-based

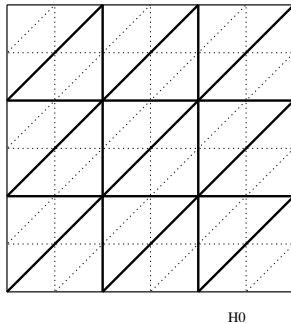
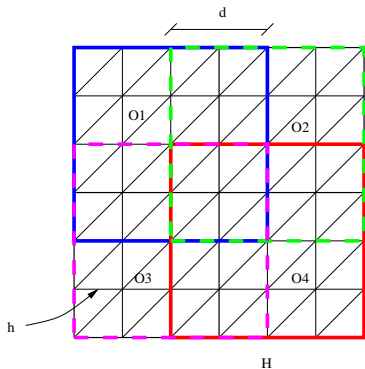




# Domain Decomposition

## Additive Schwarz methods

- Non-overlapping methods  
substructuring methods, additive average methods and others.
- Overlapping methods



## Domain Decomposition on Unstructured Grids

### Domain Decomposition on Unstructured Grids

**DOUG** (*University of Bath, University of Tartu*) 1997 - 2008

I.G.Graham, M.Haggers, R. Scheichl, L.Stals, E.Vainikko,  
M.Tehver, K. Skaburskas, O. Batrašev, C. Pöcher,

**Parallel implementation based on:**

- **MPI**
- **UMFPACK**
- **METIS**
- **BLAS**
- Automatic parallelisation and load-balancing
- Systems of PDEs
- 2D & 3D problems
- 2-level Additive Schwarz method
- 2-level partitioning
- Automatic Coarse Grid generation
- Adaptive refinement of the coarse grid



## DOUG Strategies

- Iterative solver based on Krylov subspace methods  
**PCG, MINRES, BICGSTAB**, 2-layered **FPGMRES** with left or right preconditioning.
- Non-blocking communication where at all possible  
Ax-operation:  $y := Ax - :-)$   
Dot-product:  $(x, y) - :- ($
- Preconditioner based on Domain Decomposition with 2-level solvers  
Applying the preconditioner  $P$ : solve for  $z : Pz = r . :- ?$
- Subproblems are solved with a direct, sparse multifrontal solver (UMFPACK)



## Aggregation-based Domain Decomposition methods

- Have been analysed only upto some extent
  - **We are:** making use of strong connections
  - Second aggregation for creating subdomains, or
    - using rough aggregation before graph partitioner
    - Major progress - development of the theory: **sharper bounds**
- 1 R. Scheichl and E. Vainikko, Robust Aggregation-Based Coarsening for Additive Schwarz in the Case of Highly Variable Coefficients, Proceedings of the European Conference on Computational Fluid Dynamics, ECCOMAS CFD 2006 (P. Wesseling, E. ONate, J. Periaux, Eds.), TU Delft, 2006.
  - 2 R. Scheichl and E. Vainikko, Additive Schwarz and Aggregation-Based Coarsening for Elliptic Problems with Highly Variable Coefficients, Computing, 2007, 80(4), pp 319-343.



# Aggregation

## Key issues:

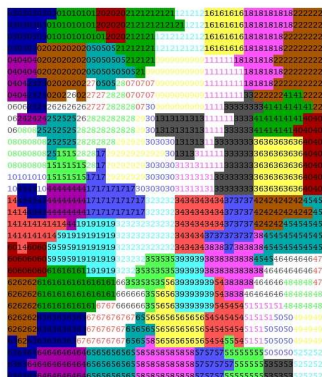
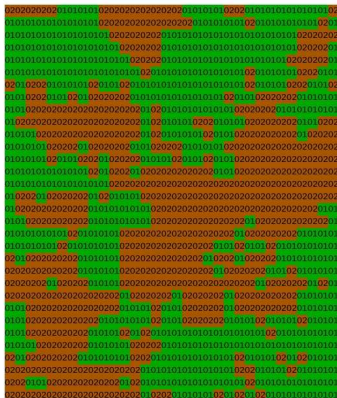
- how to find good aggregates?
- Smoothing step(s) for restriction and interpolation operators

## Four (often conflicting) aims:

- follow adequately underlying physical properties of the domain
- try to retain optimal aggregate size
- keep the shape of aggregates regular
- reduce communication => develop aggregates with smooth boundaries



# Coefficient jump resolving property



## Algorithm: Shape-preserving aggregation

**Input:** Matrix  $A$ , aggregation radius  $r$ , strong connection threshold  $\alpha$ .

**Output:** Aggregate number for each node in the domain.

1. Scale  $A$  to unit diagonal matrix (all ones on diagonal)
2. Find the set  $S$  of matrix  $A$  strong connectons:  $S = \cup_{i=1}^n S_i$ , where  $S_i \equiv \{j \neq i : |a_{ij}| \geq \alpha \max_{k \neq i} |a_{ik}|\}$ , unscale  $A$ ; `aggr_num:=0`;
3. `aggr_num:=aggr_num+1`;

Choose a seednode  $x$  from  $G$  or if  $G = \emptyset$ , choose the first nonaggregated node  $x$ ;

`level:=0`

4. If (`level<r`)

Add recursively all strongly connected non-aggregated neighbours to the aggregate `aggr_num` with `level+1` and perform smoothing step on each level

elseif (`level<2r`)

Find `layer(level+1)...layer(2r)`.

endif

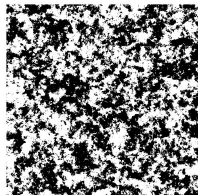
5. On the longest `layer(i)`,  $i=r+1, \dots, 2r$  add node(s) with shortest distance from  $x$  to the set  $G$  and goto step 3.



# DOUG

# Numerical results

Clipped random fields



$n = 256 \times 256$

**Time in seconds (# iterations) for different methods:**

jumpsize	Aggregation	AMG	UMFPACK
$1.5 * 10^1$	2.12 (24)	1.35 (14)	1.85
$2.2 * 10^2$	2.14 (27)	2.27 (27)	1.70
$3.3 * 10^3$	2.34 (29)	3.31 (40)	1.33
$4.9 * 10^4$	2.41 (26)	6.23 (77)	4.88

**AMG** - Aggregation-type Algebraic Multigrid [Bastian] (no smoothing - piecewise constant prolongation)

**UMFPACK** - Sparse direct solver





## DOUG on Grid

- Is it possible to use GRID in it's full power for one (huge) linear system solution?

### PROBLEM:

- Dynamic nature of GRID *versus*:
  - good parallel solvers need synchronisation steps :-(  
• no fault tolerance in mainstream MPI implementations :-(  
• => A) need for methods that do not need regular synchronisation
  - => B) Need for fault-tolerant communication libraries



## Possible solutions

### • A) Algorithms

- Asynchronous DD methods  
(based on Richardson's type iteration methods)
- Possible asynchronous Krylov subspace methods
  - Using flexible GMRES
  - some synchronisation still needed

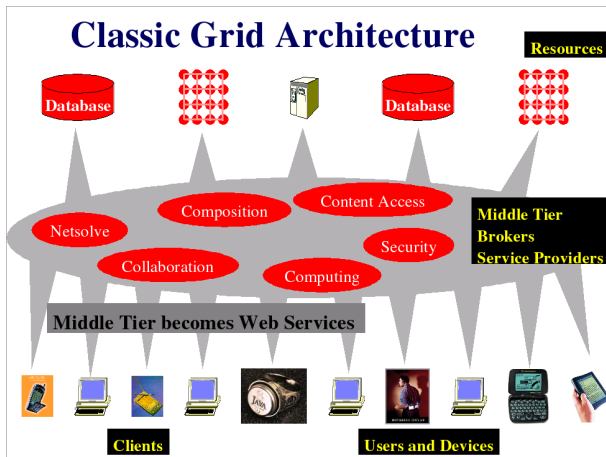
### • B) Fault tolerance

- Need for fault tolerant libraries
  - e.g. FT-MPI, OpenMPI
  - need for a standard MPI standard says currently: "*Fault tolerance is a property of an application, not the MPI Library*"
  - no FT MPI library available for GRID.



# The Grid Vision

What was promised



- As easy as power-grid
- Everywhere
- Accessible to everybody
- Access to any desired resource
  - CPU cycles
  - Storage
  - Special devices



# The Grid - what we've got so far

Batch processor

- Batch processing
  - With global scale, though
  - Parallel jobs still available, within a single cluster
  - You really need to know the resource you are running on
- Grid = merely web-services in another wrapping!
- Storage – really uncomfortable for user



## Problems with existing models

### (Not every problem in every middleware)

- Single point of failure
- Lack of scheduling
- Poor scalability
- No means for privacy
- No way for cycle stealing (scavenging)
- Firewall problem
- Requires very large installation
- No economy



# Desktop Grids

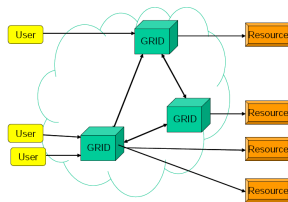
- Condor
- BOINC
  - Seti@home
  - Einstein@home
  - Proteins@home
  - Africa@home
  - SZTAKI@home
- Folding@home
- ALCHEMI
- JNGI
- **Minimum Intrusion Grid**
- **Friend-to-friend Computing**



## Minimum Intrusion Grid (MiG) Design

<http://mig-1.imada.sdu.dk/MiG/index.html>

- Non-intrusive
  - Scalable
  - Autonomous
  - Anonymous
  - Fault tolerance
  - Firewall compliant
  - Strong scheduling
    - Real-time access to CPU-s
  - Allowing multiple jobs to share a CPU for low-intensity jobs
  - Cooperative support
- BUT still, for DOUG we need (fast) communication between computing nodes!**



# Friend-to-friend Computing (F2F)

Called also: Spontaneous Desktop Grid

<http://f2f.ulno.net>

- Developed at the Institute of Computer Science, University of Tartu (U.Norbisrath and E.Vainikko)
- Computing environment based on spontaneous groups
- Instant Messaging friend groups used for triggering the Computing Grid Environment
- Merging ideas from
  - Peer-to-peer
  - High Performance Computing
  - Social networks in instant messaging
- Extremely easy concept for
  - authentication
  - authorizationbased on inherent trust between friends

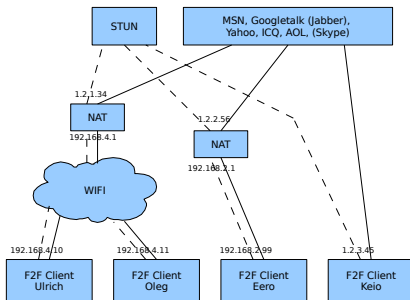




## F2F Architecture

Instant messaging protocols/solutions supported:

- MSN
- GoogleTalk (Jabber)
- Yahoo
- ICQ
- AOL
- (Skype)



# Overview

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